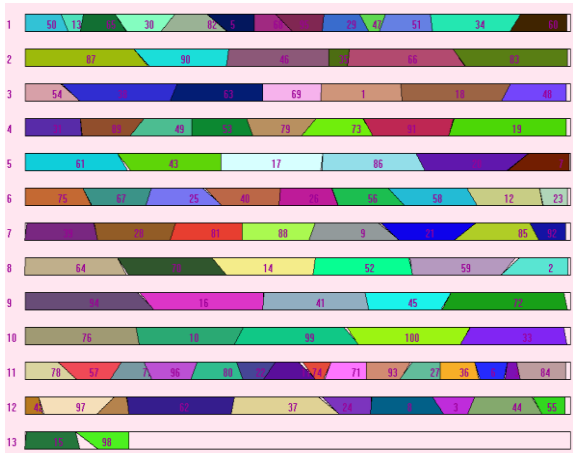
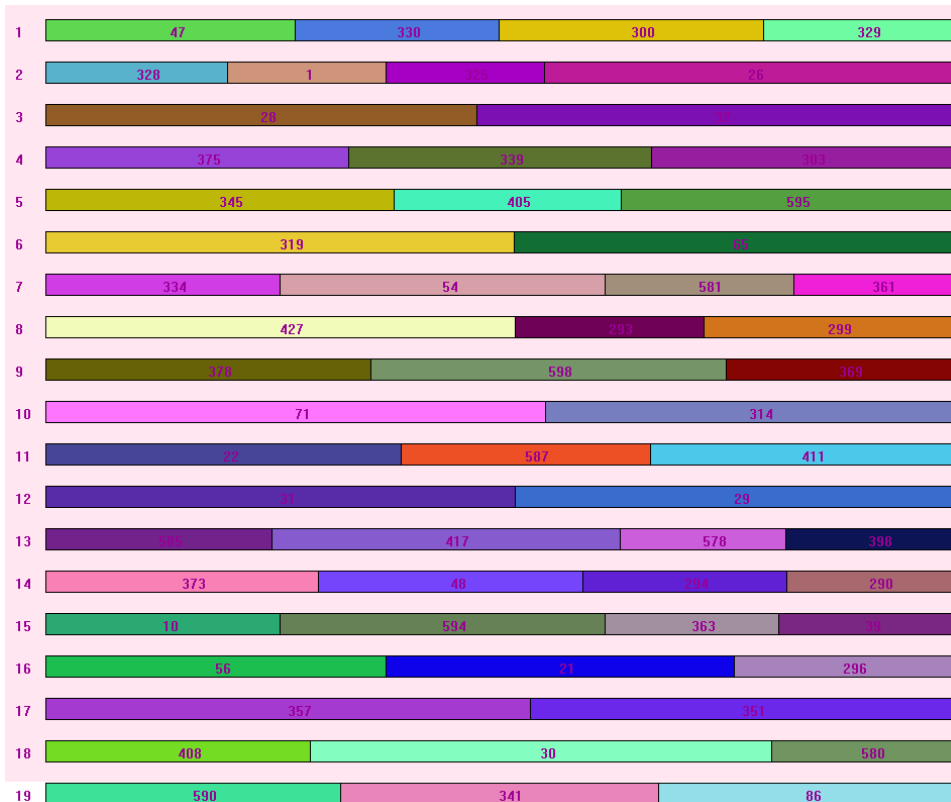


# Bin Packing with Trapezia: Methods and Applications

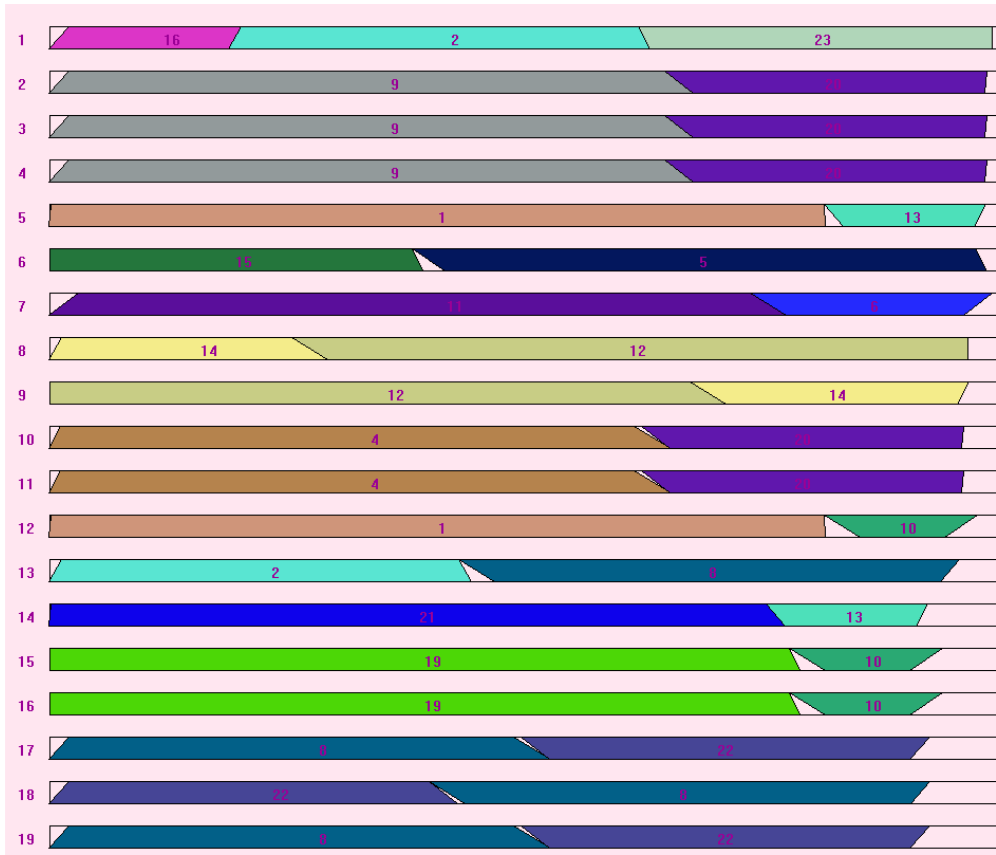


# The 1D Bin Packing Problem



- In a standard 1D bin packing problem we seek to pack 1D items into a minimum number of fixed-sized bins.
- It is NP-complete (generalises the partitioning problem).
- The order and orientation of items within each group irrelevant
- But what if it is?

# The Trapezium Packing Problem



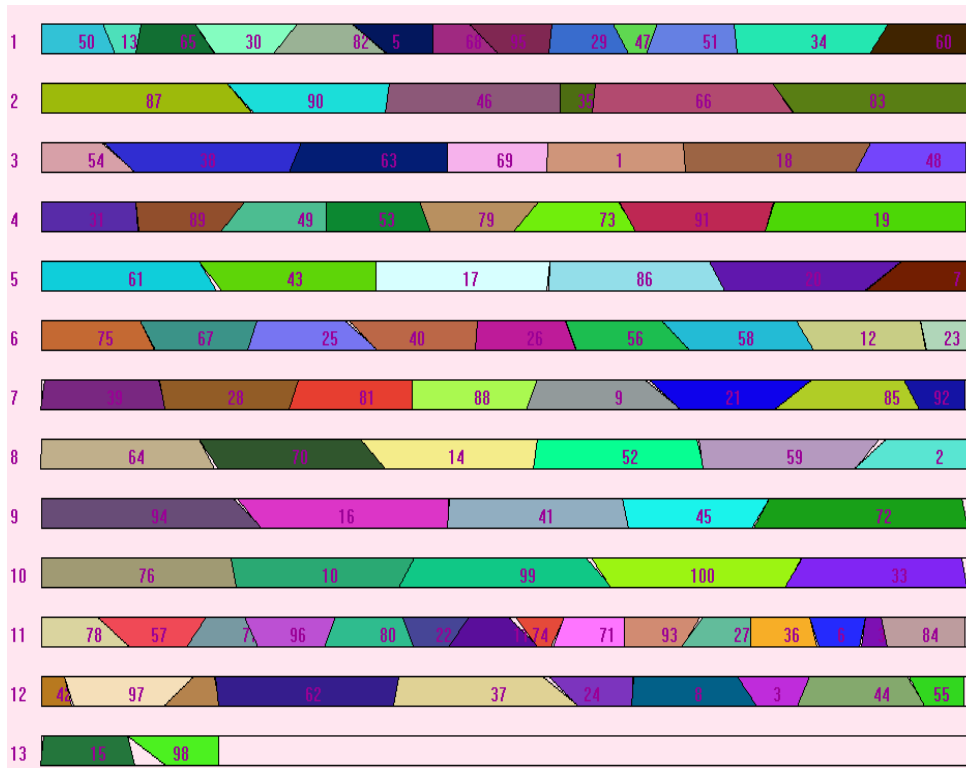
## Problem:

Arrange a set of fixed-height trapezia into a minimal number of bins

## Issues

- 1) Which items should be put into the same bin? (A bin packing problem)
- 2) How should items be arranged in a bin? (An ordering and rotation problem)

# The Trapezium Packing Problem



## Problem:

Arrange a set of fixed-height trapezia into a minimal number of bins

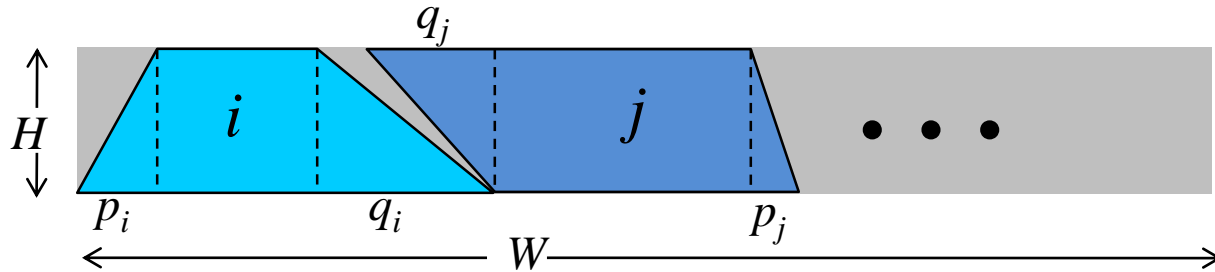
## Issues

- 1) Which items should be put into the same bin? (A bin packing problem)
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# Real World Applications



# Arranging the trapezoids in a single bin



- Items have 4 orientations, but we only need to consider two
- Inter-item wastage in the above is simply  $|q_i - q_j|$
- Wastage  $w$  = all inter-item wastage plus the LHS and RHS
- **Problem Def:** Given a group of items  $S$  such that:

Total area of trapezoids in the group

$$A(S) = \sum_{j \in S} A(j) \leq HW$$

Area of the bin

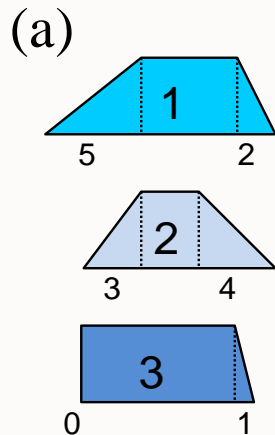
can we determine an item arrangement such that:

$$A(S) + w \leq HW$$

Total area of trapezoids in the group plus inter-item wastage

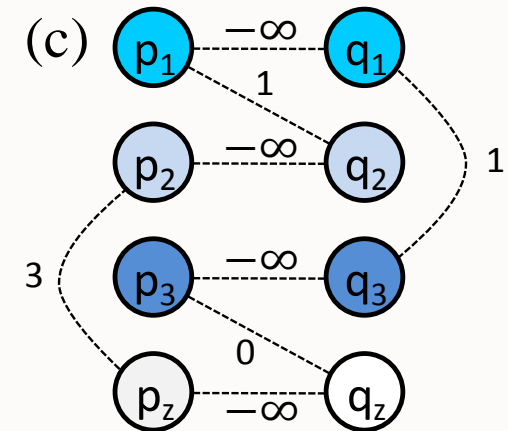
# TSP interpretation of the single bin problem

- Arranging trapezoids in a single bin can be seen as a special type of TSP
- Each projection is a “city”; edges between cities correspond to wastage; and edges on the same shape are set to minus infinity
- Can a “valid” route of less than  $w$  be achieved (discounting minus infinity arcs)?



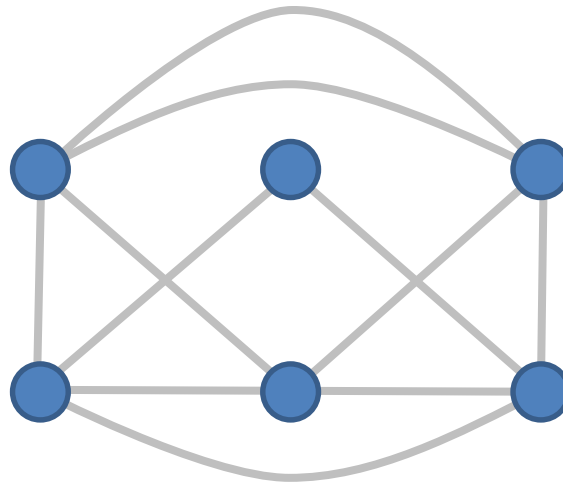
(b)

	$p_1$	$q_1$	$p_2$	$q_2$	$p_3$	$q_3$	$p_z$	$q_z$
$p_1$	$\infty$	$-\infty$	2	1	5	4	5	5
$q_1$	$-\infty$	$\infty$	1	2	2	1	2	2
$p_2$	2	1	$\infty$	$-\infty$	3	2	3	3
$q_2$	1	2	$-\infty$	$\infty$	4	3	4	4
$p_3$	5	2	3	4	$\infty$	$-\infty$	0	0
$q_3$	4	1	2	3	$-\infty$	$\infty$	1	1
$p_z$	5	2	3	4	0	1	$\infty$	$-\infty$
$q_z$	5	2	3	4	0	1	$-\infty$	$\infty$



# Revision: Eulerian Graphs and Cycles

- **Def:** An Eulerian cycle is a cycle that uses every edge in a graph exactly once.
- **Def:** An Eulerian graph is a graph that is connected and where all vertex degrees are even
- **Fact:** A graph has an Eulerian cycle if and only if it is Eulerian.





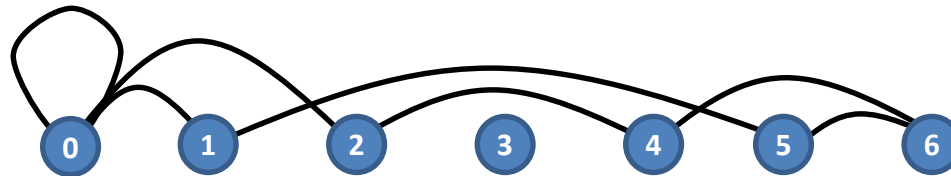
# How to Solve the Single Bin Problem

- Take a set of trapezoids



- Form a multigraph  $G$  using projection pairs as edges

$$E = \{ \{0,0\}, \{0,1\}, \{0,2\}, \{1,5\}, \{2,4\}, \{4,6\}, \{5,6\} \}$$



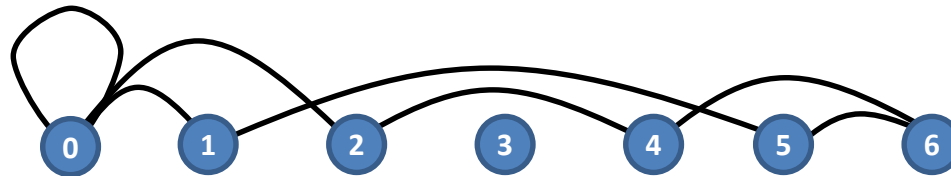
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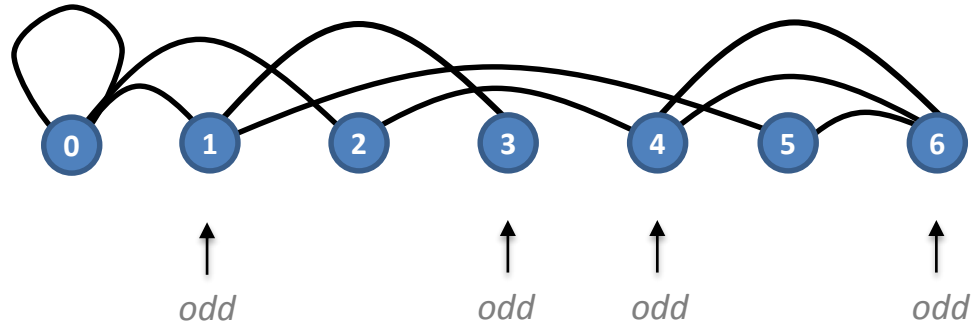


- Fact:** If  $G$  is Eulerian, then a zero cost arrangement exists.
- This is the case here:  $( (0,0), (0,1), (1,5), (5,6), (6,4), (4,2), (2,0) )$



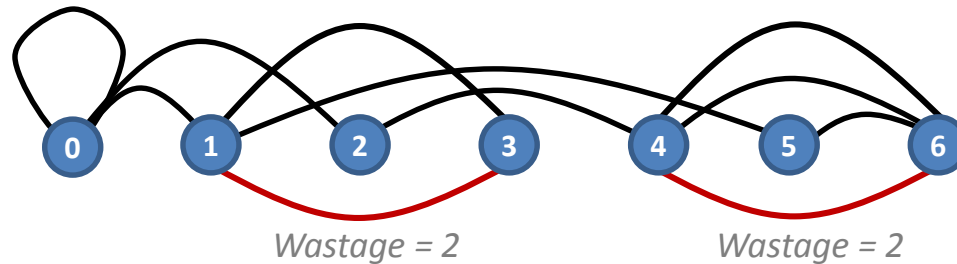
# How to Solve the Single Bin Problem

- **Fact.** If  $G$  has odd-degree vertices, then a minimum weight matching can be introduced to make them even.
- For example,



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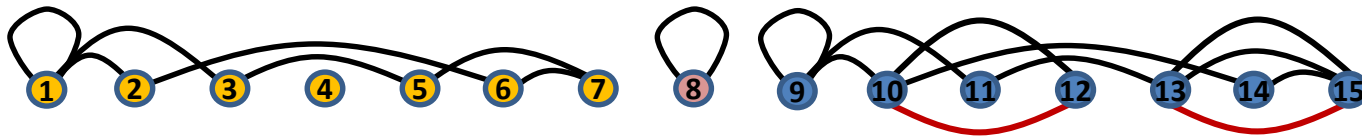
- If  $G$  is now connected, an optimal arrangement is an Eulerian cycle in this new graph as before

$((0,0),(0,2),(2,4),(4,6),(6,4),(6,5),(5,1),(1,3),(1,0))$

$\underbrace{\hspace{10em}}$   
Wastage = 2
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# How to Solve the Single Bin Problem

- If all degrees are even but  $G$  is disconnected, the components need to be joined
- **Fact.** Each of the  $n$  Eulerian components is optimal.
- **Fact.** Optimally joining two components gives us  $n - 1$  optimal Eulerian components



$w = 4$

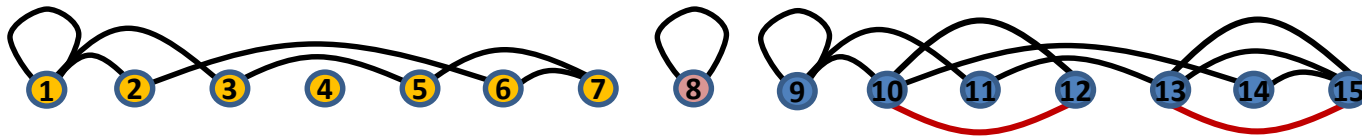
•  $((1,1),(1,3),(3,5),(5,7),(7,6),(6,2),(2,1))$  *Wastage = 0*

•  $((8,8))$  *Wastage = 0*

•  $((9,9),(9,11),(11,13),(13,15),(15,13),(15,14),(14,10),(10,12),(10,9))$  *Wastage = 4*

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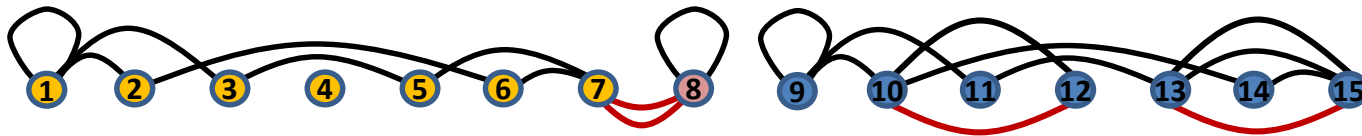
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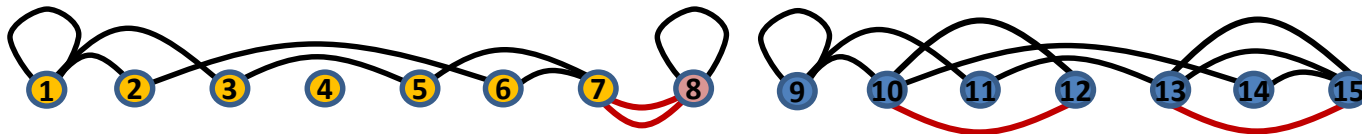
$w = 6$

- $((1,1),(1,3),(3,5),(5,7), (8,8), (7,6),(6,2),(2,1))$  *Wastage = 2*

- $((9,9),(9,11),(11,13),(13,15),(15,13),(15,14),(14,10),(10,12),(10,9))$  *Wastage = 4*

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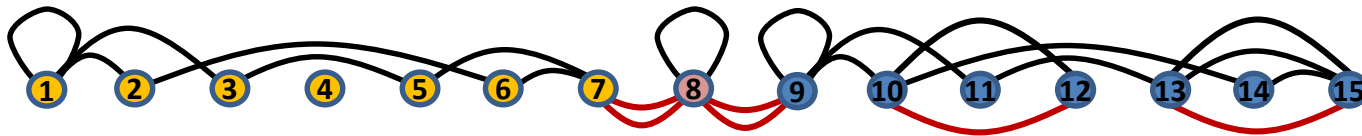
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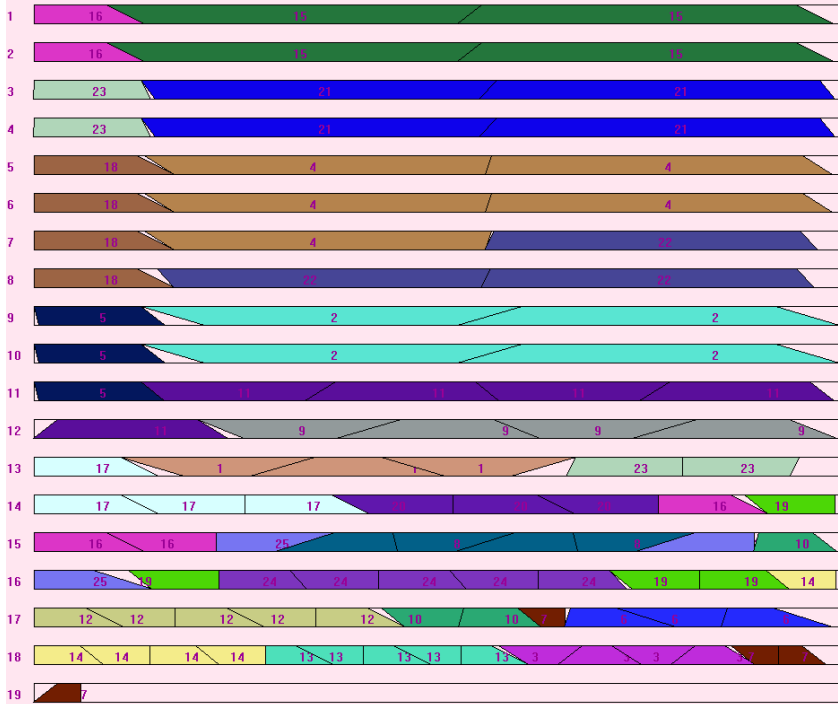


$w = 8$

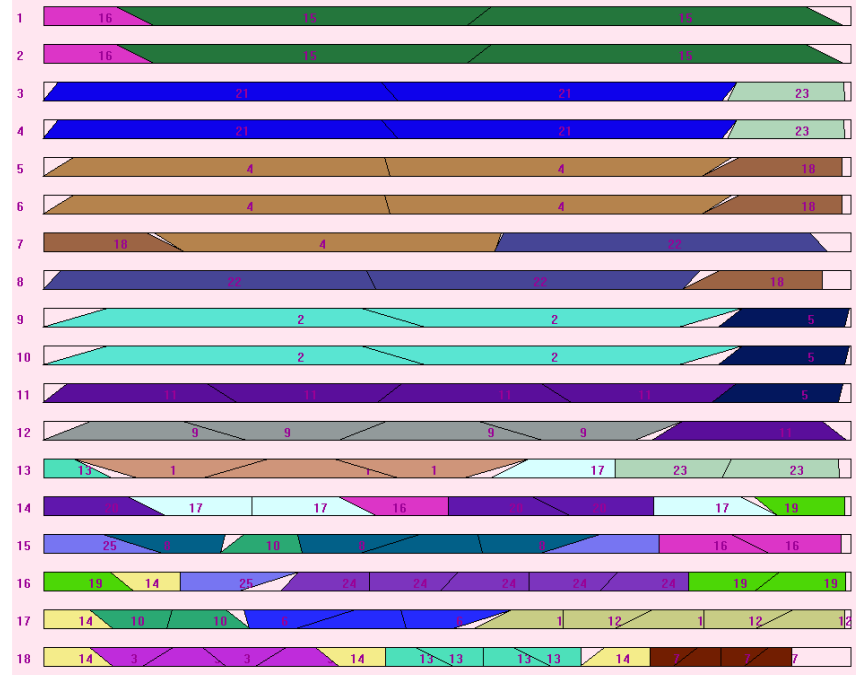
- $((1,1),(1,3),(3,5),(5,7), (8,8) (9,9),(9,11),(11,13),(13,15),(15,13), (15,14),(14,10),(10,12),(10,9) ,(7,6),(6,2),(2,1))$  *Wastage = 8*

- A proof naturally follows which demonstrates this method to be exact. The problem using multiple bins is still NP-hard, however.

# The multi-bin problem: Comparison using FFD



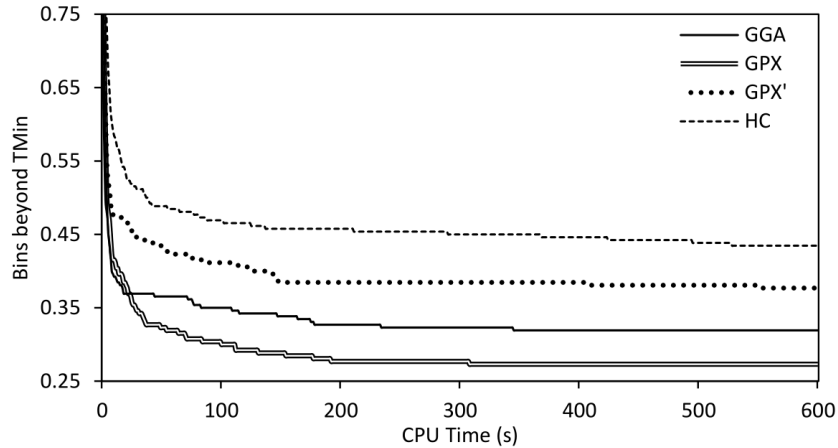
- Example solution using a simple (inexact) heuristic to pack trapezia into individual bins, combined with FFD.



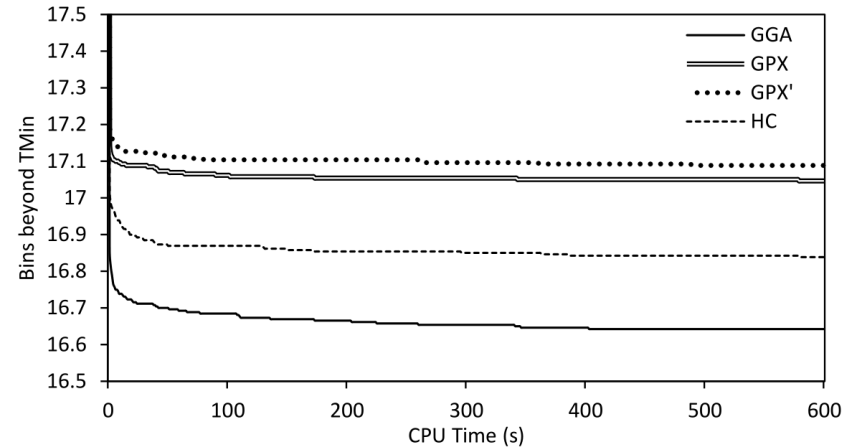
- Solution using the exact algorithm for packing trapezia into individual bins, combined with FFD

# Further Improvements via Evolutionary Methods and Local Search

Run profiles for various different evolutionary operators using 500-item problems (averaged across 240 instances)



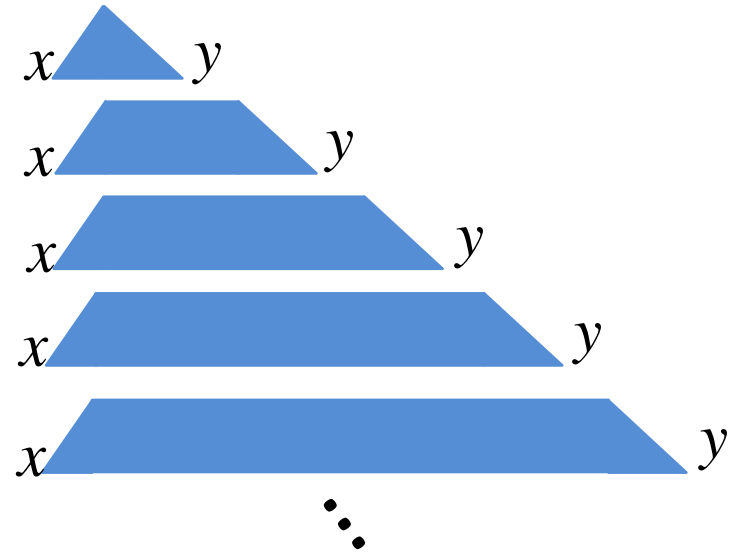
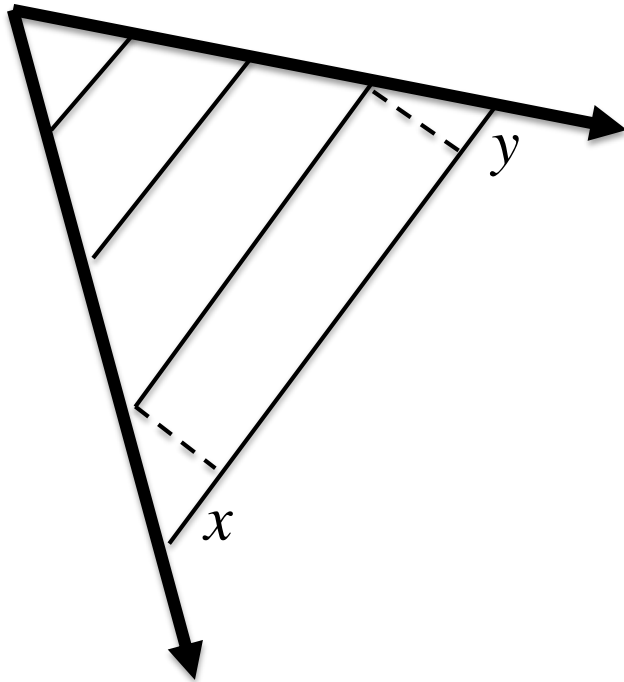
Instances with approximately  
9 items per bin



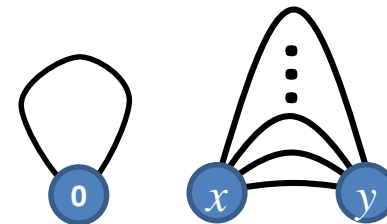
Instances with approximately  
2.5 items per bin

- **Lewis, R.** and P. Holborn (2017) 'How to Pack Trapezoids: Exact and Evolutionary Algorithms'. *IEEE Transactions on Evolutionary Computation*, vol. 21(3), pp. 463-476.

# Special Cases Arising in Industry

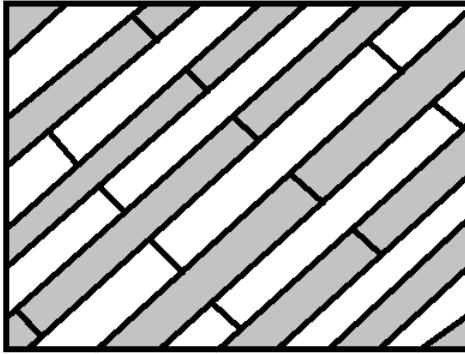
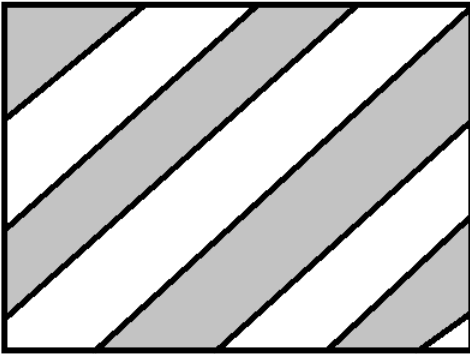
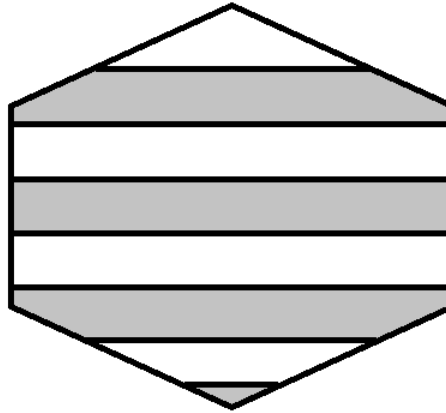
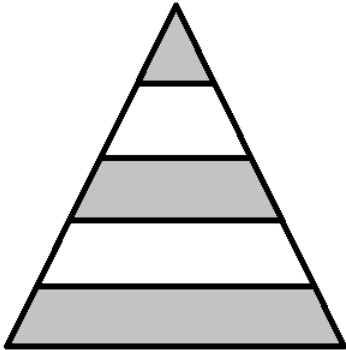


If item sizes in a 1D BPP are  $1, 2, 3, \dots, j$ , then a perfect packing exists if and only if the sum of all item sizes is a multiple of the bin capacity



- Coffman, E. et al. "Perfect Packing Theorems and the Average-Case Behaviour of Optimal and Online Bin Packing" SIAM Review 2002 44:1, 95-108

# Special Cases Arising in Industry



# Bin Packing with Trapezia: Methods and Applications

Questions...

